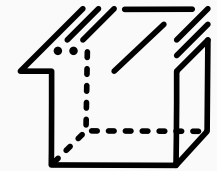


Towards Computational UIP in Cubical Agda



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DistriNet, KU Leuven, Belgium

Cubical Agda [VMA21]

Based on [Coh+17] with De Morgan cubes.

- `transp` and `hcomp`
- *I*, *Glue* types, higher inductive types

Glue types (source of univalence) incompatible with Uniqueness of Identity Proofs (UIP) and K.

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Glue types (source of univalence) incompatible with Uniqueness of Identity Proofs (UIP) and K.
(But we can disable it!)

FunExt, **Quotients**, possibly **UIP**... good for working with sets! [Coc19, Shu17]

Motivation

Choices for working with **sets in Agda** (with Uniqueness of Identity Proofs):

	FunExt	UIP holds	UIP computes	Quotient	Projects
Agda --with-K					(Many)
Agda + Setoids		Setoid 		Setoid 	[CND25, HC21]

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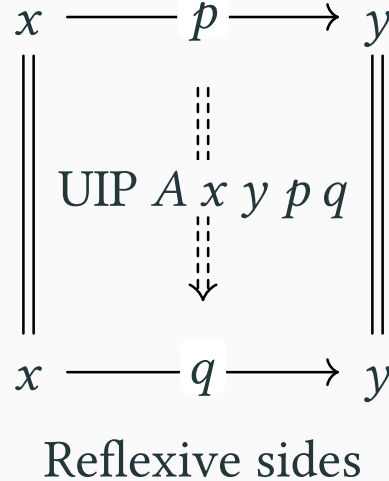
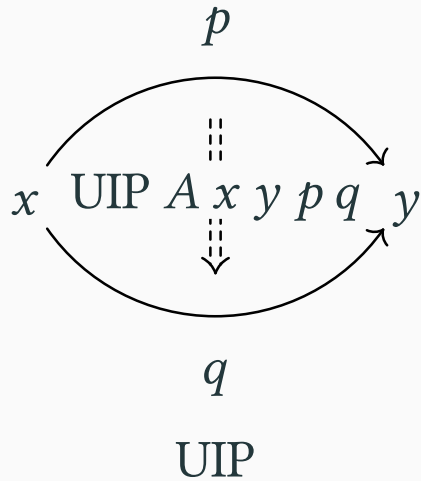
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Cubical-UIP

- ✓ Agda `--cubical=no-glue` variant (merged in 2.9.0)
- 🕒 Agda `--cubical=uii` variant with a “UIP primitive” that recurses on types

UIP: working with h -Sets

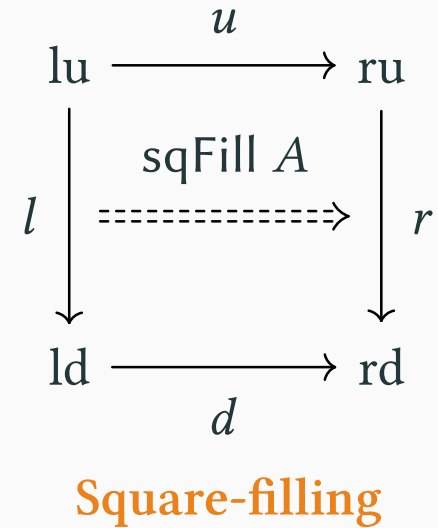
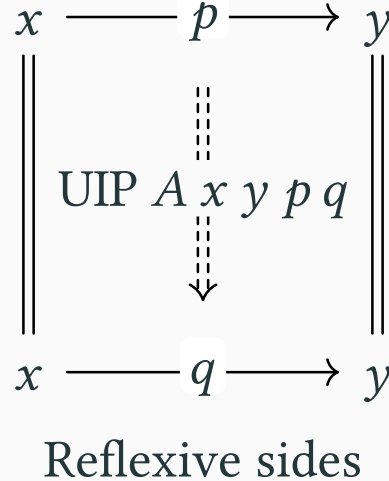
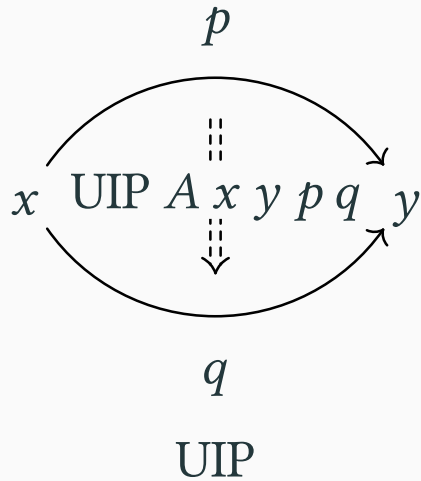
UIP: squares with reflexive sides have a filler.



¹[Agd25]

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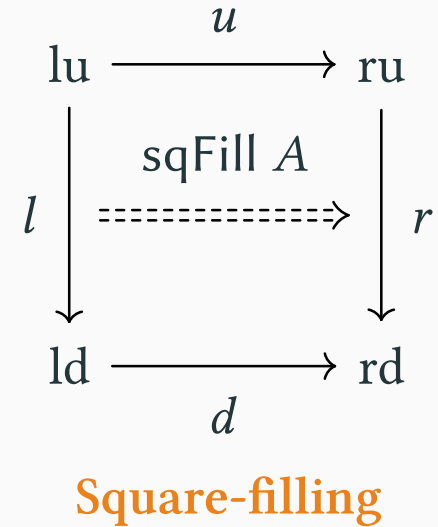
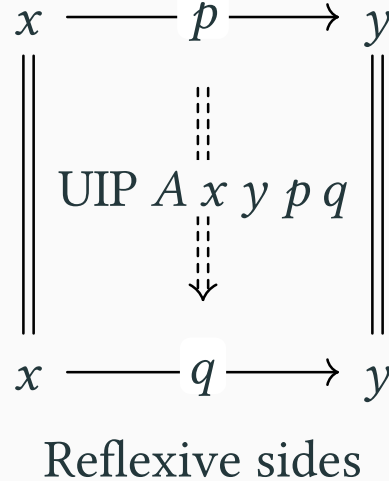
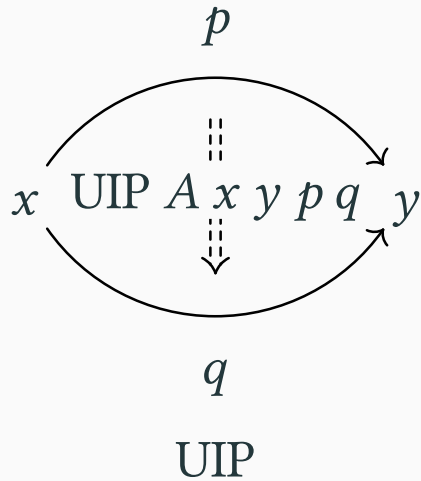


SqFill: any given square has a filler.

¹[Agd25]

UIP: working with h -Sets

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SqFill: any given square has a filler.

UIP \leftrightarrow ¹ SqFill

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Implemented: Pi, Sigma, PathP; List, Maybe, (Co-)products ; Base types: Unit, Bool, Nat

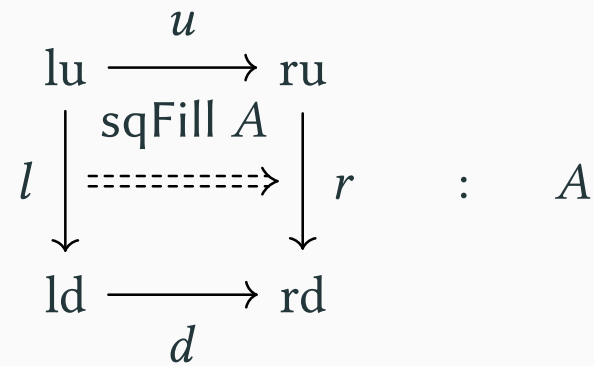
Welcome to experiment: <https://github.com/yeejian-tan/agda/tree/cubical-uip>

Non-dependent vs Dependent

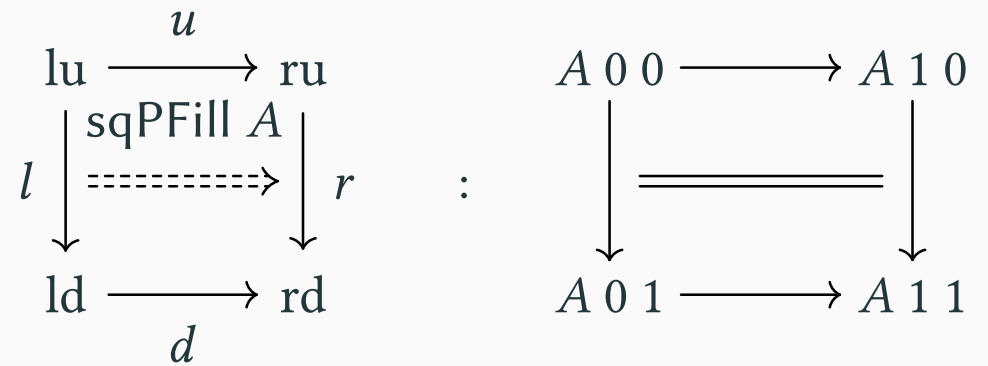
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Homogenous SqFill

Non-dependent vs Dependent

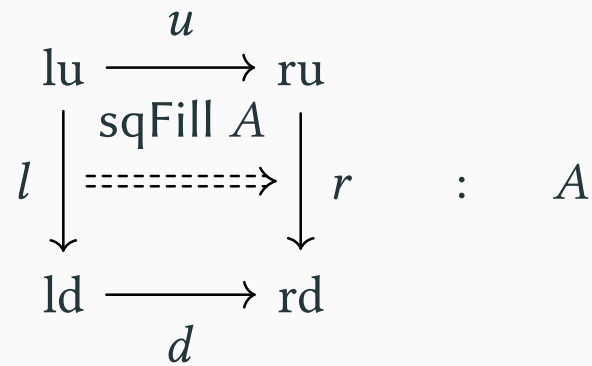


Homogenous SqFill

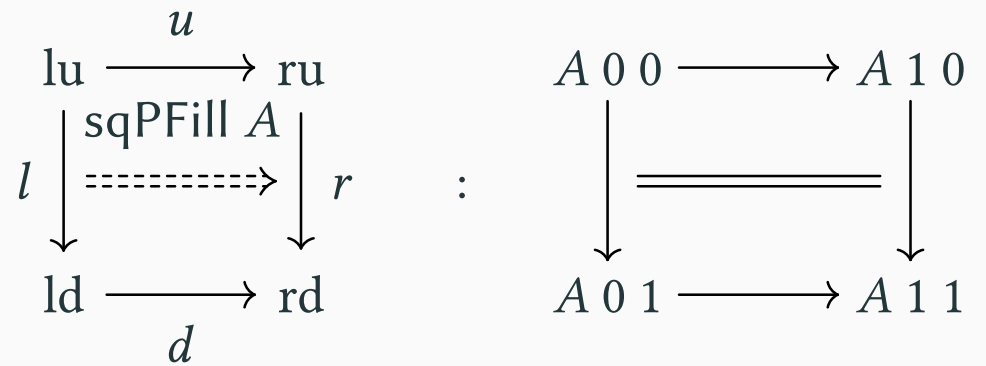


Heterogeneous SqPFill

Non-dependent vs Dependent



Homogenous SqFill



Heterogeneous SqPFill

SqPFill (heterogenous) \leftrightarrow pointwise SqFill

Observation: Complexity of Preservation Proofs

Our proofs attempt to fill squares directly instead of instantiating more general results [TND25].

Preservation proofs differ in complexity for SqFill and SqPFill:

	SqFill (non-dependent)	SqPFill (dependent)
Pi	Trivial (no Kan operations)	Complicated: transport-fill-align
Sigma	Complicated: transport-fill-align	Trivial (no Kan operations)
Coproducts	Encode-decode (J , hcomp) [Uni13]	Encode-decode (J , comp) [Uni13]
Path Types	Simple (a single hcomp)	Complicated: transport-fill-align

Complexity of SqFill and SqPFill preservation proofs

Candidates:

- Chen, Nordvall Forsberg, and Tsai's TT-as-QIIRT [CNT26]
Replacing UIP postulate by our sqFill primitive
- Bitvector Library in Katamaran [Huy+23]
computation over proofs of “proof-irrelevance”

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Do you have evaluation targets that **compute** over **UIP proofs**?

- Possibly need: **FunExt**, **Quotient types**?
- Would benefit from **computational UIP/proof-irrelevance**?
- No need: **univalence**?

Related Work

System	UIP	Equality ... compute
Cubical-UIP	Propositional (+ automation)	proofs

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XTT	Definitional	proofs

XTT [SAG22]

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- type-checking algorithm requires inj. of type formers (contradicts at least LEM)

System	UIP	Equality ... compute
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XTT	Definitional	proofs
Observational Type Theory (OTT)	Definitional (needs SProp)	types

XTT [SAG22]

- *Definitional* UIP (c.f. our propositional UIP)
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Observational Type Theory [Alt99, AMS07, PLT25]

- UIP needs proof-irrelevant propositions
- Equality **types** compute instead of **proofs**

$$\sim_{A \times B} = \sim_A \times \sim_B$$

(c.f. *isomorphic* in Cubical)

What happens to $\text{sqFill } U$?

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Option 1: U is an inductive-recursive universe

```
data U : Type1 where
  Pi : (A : U) → (B : El A → Type) → U
  ...
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Advantage: `sqFill U` (and `hcomp U`) compute by recursing on typecodes.

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- `hcomp U` does *not* compute in other Cubical Agda variants
- anti-classical: contradicts with at least excluded middle.

Open question: Universe

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- | | |
|--|--------------------------|
| | no-glugue (?) |
| | / \ |
| | full (TF) uip (computes) |

Option 2: Universe is *not* closed

- `sqFill U` as a new type... ?
- `sqFill` only for small types?

Future work

Universes.

Records: nested Σ -types?

Datatypes:

- HITs: add sqFill constructor
- Non-HITs: add sqFill constructor; or “pattern-match” on corners (like hcomp)

Contribution

 `--cubical=no-glue`: (Cubical Agda - Glue)

 `--cubical=uiop`: (Cubical Agda - Glue + sqFill primitive recursing on types) [TND25]

Supports (Bool, Nat, Unit), (Pi, Sigma, PathP, List, Maybe, (Co-)products).

Welcome to experiment: <https://github.com/yeejian-tan/agda/tree/cubical-uiop>

Questions?

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See a square compute! 👁👁